

## NEIGHBORHOODS AND COVERING VERTICES BY CYCLES

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In this work, we show the following results. First if  $G = (V, E)$  is a 2-connected graph, and  $X$  is a set of vertices of  $G$  such that for every pair  $x, x'$  in  $X$ ,  $|N_G(x) \cup N_G(x')| \geq n/2 + 2$ , and the minimum degree of the induced graph  $\langle X \rangle$  is at least 3, then  $X$  is covered by one cycle.

This result will be in fact generalised by considering tuples instead of pairs of vertices.

Let  $\delta_1(X)$  be the minimum degree in the induced graph  $\langle X \rangle$ . For any  $t \geq 2$ ,

$\delta_t(X) = \min\{|N_G(u_1) \cup N_G(u_2) \dots \cup N_G(u_t)|, u_i \neq u_j, u_1, \dots, u_t \in X\}$ .

If  $\delta_1(X) \geq t$ , and  $\delta_t(X) \geq |V|/p + t$ , then  $X$  is covered by at most  $(p-1)$  cycles of  $G$ . If furthermore  $\delta_1(X) \geq 2t$ ,  $(p-1)$  cycles are sufficient.

So we deduce the following:

Let  $p$  and  $t$  ( $t \geq 2$ ) be two integers.

Let  $G$  be a 2-connected graph of order  $n$ , of minimum degree at least  $t$ . If  $\delta_1 \geq t$ , and  $\delta_t \geq n/p + t$ , then  $V$  is covered by at most  $(p-1) + \lceil (t-1)/k \rceil$  cycles, where  $k$  is the connectivity of  $G$ .

If furthermore  $\delta_1 \geq 2t$ ,  $(p-1)$  cycles are sufficient.

In particular, if  $\delta_1 \geq 2t$  and  $\delta_t \geq n/2 + t$ , then  $G$  is hamiltonian.

**1. Introduction**

In this work we consider finite and undirected simple graphs. If  $G = (V, E)$  is a graph we denote  $N_G(z)$  the neighborhood of a vertex  $z$  i.e. the set of vertices of  $G$  adjacent to  $z$ .

Let  $N_G[a] = N_G(a) \cup \{a\}$ ,  $N_G(a \cup b) = N_G(a) \cup N_G(b)$ , and  $\delta_2 = \min\{|N_G(u) \cup N_G(v)|, u, v \in V \text{ and } u \neq v\}$ .

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For any subset  $X$  of  $V$ , then  $\langle X \rangle$  is the subgraph of  $G$  induced by the set  $X$ ,  $\delta_1(X)$  is the minimum degree in the graph  $\langle X \rangle$ . For any  $t \geq 2$ ,  $\delta_t(X) = \min\{|N_G(u_1) \cup N_G(u_2) \dots \cup N_G(u_t)|, u_i \neq u_j, u_1, \dots, u_t \in X\}$ , and

$$\sigma_t(X) = \min\left\{\sum_{x \in S} \deg_G x \mid S \subseteq X \text{ is an independent set in } G \text{ and } |S| = t\right\}.$$

Note that  $\sigma_t(X) \geq \delta_t(X)$ . If  $X = V$ , we will denote  $\delta_t(X)$  by simply  $\delta_t$ .

A  $p$ -cycle cover of  $X \subset V$  in the graph  $G$  is a family of  $p$  cycles of the graph  $G$  such that each element of  $X$  belongs to at least one of the cycles of this family. We say also that  $X$  is covered by  $p$  cycles.

The maximum length of a cycle in  $G$  is called circumference of  $G$ .

Let  $C$  be a cycle. A path  $Q[a, b]$  of endpoints  $a$  and  $b$  is said to be strongly joined to the cycle  $C$  if  $Q$  has no vertex in common with  $C$  and there exist 2 vertex-disjoint paths  $P[a, c]$  and  $P[b, c']$  from  $Q[a, b]$  to  $C$  internally disjoint from  $C$ . These two last paths are internally vertex disjoint from  $Q[a, b]$ . Then  $Q[a, b]$  is called a  $C$ -path and the vertices  $c$  and  $c'$  are called joins of the  $C$ -path  $Q[a, b]$ .

We suppose that  $C$  has an orientation and a vertex labeling following this orientation. We denote by  $[a, b]$  (resp.  $[a, b[$ ) the segment of  $C$  consisting of the vertices  $x$  such that  $a \leq x \leq b$  (resp.  $a \leq x < b$ ). We denote by  $]a, b[$  the segment  $[a, b] \setminus \{a, b\}$ . Two chords  $au$  and  $bu'$  of  $C$  are crossings if we meet successively  $(a, u', u, b)$  or  $(a, b, u, u')$ .

We recall the definition of a quasi-claw free graph  $H$ . For each pair of vertices  $a, b$  of  $H$  at distance 2, there exists at least a vertex  $u \in N[a] \cap N[b]$  such that  $N[u] \subset N[a] \cup N[b]$ .

Several authors studied relations between the parameter  $\delta_2$  and hamiltonicity, or the circumference of the graph. It has been shown that

**Theorem A.** [3] *Let  $G$  be a 2-connected graph of order  $n$ . If  $\delta_2 \geq n/2$ , then  $G$  is hamiltonian for  $n$  sufficiently large.*

Jackson [4] has proved that a 3-connected graph with  $\sigma_2 \geq (n+1)/2$  is hamiltonian.

**Theorem B.** [2] *If  $G$  is a 2-connected graph of order  $n$ ,  $n \geq 10$ , and  $\delta_2 \geq 2n/5 + 1$ , then there exists a 2-cycle cover of  $G$ . Furthermore, one of the cycles can be chosen as a longest cycle of  $G$ .*

On the other hand, we have also the result [5]:

“If  $\sigma_t(G) \geq n$ , then  $G$  is covered by at most  $(t-1)$  subgraphs, each of them is a cycle, an edge or a vertex”. More exactly,

**Theorem C.** [5] *Let  $k \geq 2$  be an integer. Let  $G$  be a graph on  $n$  vertices and let  $X \subseteq V$ . If  $\sigma_t(X) \geq n$  or  $\alpha(X) < t$  then  $X$  is covered with  $t-1$  cycles, edges or vertices of  $G$ .*

## 2. Main results

First we obtain the following results for 2-connected graphs.

**Theorem 1.** *Let  $p$  ( $p \geq 2$ ) be any integer. Let  $G$  be a 2-connected graph of order  $n$ , and of minimum degree 3. If  $\delta_2 \geq n/p+2$ , then the vertices of  $G$  are covered by at most  $p-1$  cycles.*

This theorem is sharp. Let us remark that the Petersen graph satisfies  $\delta_2 = n/2$  but is not hamiltonian.

For  $p = 2$  and 3, this theorem gives an improvement of [Theorems B and C](#). In the particular case where  $G$  is a quasi-claw-free graph, this result recalls a result of [1]. In fact, with the hypotheses of the theorem, we have  $\sigma_{2p}(G) \geq n$ . In a set of  $2p$  vertices of  $G$ , either there exist 2 adjacent vertices or 2 vertices at distance 2. So  $\alpha(G^2) \leq 2p-1$ . Then by [1], as  $G$  is 2 connected,  $V$  is covered by at most  $p$  cycles.

To establish [Theorem 1](#), we prove the following basic result:

**Proposition 1.** *Let  $G = (V, E)$  be a 2-connected graph of order  $n$ , and  $X$  a set of vertices  $G$ . Suppose that, for every pair  $x, x'$  in  $X$ ,  $|N_G(x) \cup N_G(x')| \geq n/2 + 2$  and the minimum degree of  $\langle X \rangle$  is at least 3. Then  $X$  is covered by one cycle of  $G$ .*

We generalize [Theorem 1](#) as follows.

**Theorem 2.** *Let  $p$  and  $t$  ( $p \geq 2, t \geq 3$ ) be two integers. Let  $G$  be a 2-connected graph of order  $n$ , such that  $\delta_t \geq n/p+t$ .*

1) *If the minimum degree of  $G$  is at least  $t$ , then the vertices of  $G$  are covered by at most  $(p-1) + \lceil (t-1)/k \rceil$  cycles, where  $k$  is the connectivity of  $G$ .*

2) *If, furthermore, the minimum degree of  $G$  is at least  $2t$ , then the vertices of  $G$  are covered by at most  $(p-1)$  cycles.*

This theorem is sharp. Let  $k$  ( $k \geq 2$ ) be an integer. Let  $T$  be a set of  $k$  independent edges and let  $S$  be an independent set of  $(k-1)$  other vertices. Let  $G_0$  be the join of  $S$  and  $T$ . The graph  $G_0$  is of minimum degree  $k$ . Then  $\delta_k = (2k - 1) = k + (k-1)$ , this value corresponds to the order of the neighborhood of  $k$  vertices of  $T$ . We have  $n = 3k-1$  and  $p = 3$ . The theorem gives a 2-cycle cover of  $G_0$ , and this is in fact a minimum covering.

We prove [Theorem 2](#) in the following form.

**Proposition 2.** Let  $p \geq 2$  and  $t \geq 3$  be two integers. Let  $G$  be a 2-connected graph of order  $n$ , and  $X$  a set of vertices of  $G$ .

1) If the minimum degree in the subgraph  $\langle X \rangle$  is at least  $t$ , and for each tuple in  $X$ , we have  $|N_G(x_1) \cup N_G(x_2) \dots \cup N_G(x_t)| \geq n/p + t$  then

$X$  is covered by at most  $(p-1) + \lceil (t-1)/k \rceil$  cycles, where  $k$  is the connectivity of  $G$ .

2) If furthermore, the minimum degree of  $\langle X \rangle$  is at least  $2t$ , then  $X$  is covered by at most  $(p-1)$  cycles.

### 3. Proofs

We give now the proof of Proposition 1.

**Proof of Proposition 1.** The proof is by contradiction.

Let  $C$  be a cycle of  $G$

i) containing the maximum number of vertices of  $X$ ,

ii) and then,  $C$  is, among the cycles satisfying (i), of maximum length.

As  $C$  does not cover  $X$ , by the 2-connectivity of  $G$ , there exists a  $C$  path intersecting  $X \setminus C$ . We give an orientation to this cycle.

*First Case.*  $X \setminus C$  is not an independent set.

So  $X \setminus C$  contains at least 2 vertices. Let  $Q$  be a path such that

$\alpha$ )  $Q$  is strongly joined to  $C$ .

$\beta$ )  $|Q \cap X|$  is maximum,

Let  $x_1x_2$  be an edge of  $\langle X \setminus C \rangle$ . Then the path  $x_1x_2$  is strongly joined to the cycle  $C$ . So, by maximality of  $Q$ , we have  $|Q \cap X| \geq 2$ . We may suppose  $Q = Q[x, x']$  with  $x$  and  $x'$  vertices of  $X$ . Let  $c, c'$  be the joins of  $Q$  with  $C$  chosen so that there is no neighbor  $c''$  of  $x$  or  $x'$  in  $]c, c'[$  ( $\gamma$ ).

By maximality of  $C$ , each of the segments  $]c, c'[$  and  $]c', c[$  intersect  $X$ .

Let  $a$  (respectively  $a'$ ) be the first vertex of  $X$  which follows  $c$  (respectively  $c'$ ) on the cycle.

And also, by maximality of  $C$  and by minimality (see  $\gamma$ ) of  $]c, c'[$ , we have

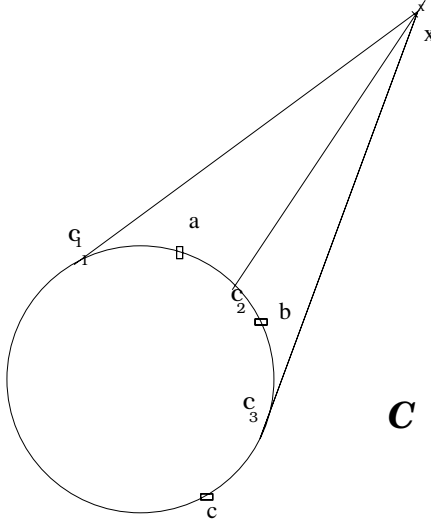
$$N_C(x \cup x') \cap N_C(a)^- \subset \{c, c'\} \quad (1)$$

$$N_C(x \cup x') \cap N_C(a')^- \subset \{c, c'\} \quad (2).$$

Similarly, there is no neighbor contained in  $G \setminus C$  common to some of the vertices  $x, x'$  and to some of  $a, a'$ . By hypothesis on  $X$ , we have

$$(*) \quad |N_G(x \cup x')| \geq n/2 + 2$$

$$(**) \quad |(N_{G \setminus C}(a \cup a'))| + |N_C(a \cup a')^-| \geq n/2 + 2.$$



**Fig. 1.** The cycle  $C$

By maximality of  $C$ , we observe that  $N_C(x \cup x')$  does not contain  $a$  or  $a'$ ; The set  $A = (N_{G \setminus C}(a \cup a') \cup N_C(a \cup a'))^-$  does not contain  $x$  or  $x'$ . So from  $(*)$  and  $(**)$ , it follows that each of the sets  $A$  and  $N_G(x \cup x')$  has at least  $n/2$  vertices in  $G \setminus \{a, a', x, x'\}$ , they have an intersection of at least 4 vertices. We get a contradiction with the inequalities (1) and (2).

Case 1 is now excluded. Let  $d(C) = \min\{d_C(c, c'), c \text{ and } c' \text{ pair of joins}\}$ . If there are several choices, we choose  $C$  such that  $d(C)$  is minimal.

*Second Case.*  $X \setminus C$  is an independent set.

Let  $x$  be a vertex in  $X \setminus C$ . By hypothesis on  $X$ , the vertex  $x$  has at least three neighbors  $c_1, c_2, c_3$  in  $X$ , then on the cycle  $C$ . We choose  $x$  such that the distance between two neighbors (say  $c_2, c_3$ ) is minimum. Let  $a, b, c$  be the vertices of  $X$  which follow respectively  $c_1, c_2, c_3$  on  $C$ .

Let  $A = (N_{[c_1^{++}, b^-]}(b \cup c))^+$ .  $B = (N_{[c^+, c_1^+]}(b \cup c))^+$ ,  $C = (N_{[b^+, c]}(b))^-$ ,  $D = (N_{[b, c^+]}(c))^+$ .

Then,

$\alpha$ ) by maximality of  $C$ ,  $b$  has no neighbor in  $]c_3, c[$ , neither in  $]c_1, a[$ .

$\beta$ ) by definition of  $C$  there is no triple  $u^-, u, u^+$  in  $[b, c_3[$  such that  $u^-$  is neighbor of  $c$ , and  $u^+$  is neighbor of  $b$ . Otherwise, either we get a longer cycle, or, a cycle  $C_1$  with the same length as  $C$  and such that  $d(C_1) < d(b, c_3) < d(C)$ ; we are in contradiction with the definition of  $C$ . We remark that  $bc_3^+$  is not an edge. So, in the segment  $[b, c^-]$ , there is no common vertex to  $(N(b))^-$  and  $(N(c))^+$ .

Thus  $C \cap D$  is empty. It is obvious that  $A \cap B$  is empty.

Consequently  $|A \cup B| \geq |N_{[c^{++}, b^-]}(b \cup c) \cup N_{[c^+, c_1^+]}(b \cup c)|$ .

$|C \cup D| \geq |N_{[b^+, c]}(b) \cup N_{[b, c^-]}(c)|$

By maximality of  $C$ , we get  $(A \cup B) \cap (C \cup D) \subset \{b \cup c\}$ .

Using the hypothesis on the neighborhoods of  $X$ , we conclude that the set  $M = A \cup B \cup C \cup D \cup N_{G \setminus C}(c \cup b)$  has at least  $n/2$  vertices. (3)

On the other hand, by maximality of the cycle  $C$ , there is no crossing between 2 chords of  $C$ , of the form  $au \ bu'$  (where  $u' = u^-$  or  $u^+$ ). Thus

$$N_C(a) \cap (N_{[b^+, c_1^+]}(b))^- = \emptyset.$$

$$N_C(a) \cap (N_{[c_1^{++}, (b)^-]}(b))^+ = \emptyset.$$

Analogously the same holds when we replace  $b$  by  $c$ .

If we replace  $a$  by  $x$  in the last two sets, the only possible intersection is in  $\{c_2, c_3\}$ . Then

$$N_C(x \cup a) \cap (A \cup B \cup C \cup D) \subset \{c_2, c_3\}. \quad (4)$$

$$\text{Furthermore, } N_{G-C}(x \cup a) \cap N_{G-C}(c \cup b) = \emptyset \quad (4')$$

By hypothesis on  $X$ ,  $|N_G(x \cup a)| \geq n/2 + 2$  vertices. The two sets  $M$  and  $N_G(x \cup a)$  are subsets of  $G \setminus \{x\}$ . So, by (3), the intersection of these 2 sets must contain at least 3 vertices.

We get also a contradiction with (4) and (4') •

### Remark

1) The proposition is sharp. For example in the graph  $K_{(p, p+1)}$ , we cannot cover the stable set  $X$  of cardinality  $p+1$  by one cycle; and  $\delta_2 \geq (n-1)/2$ .

2) We can replace the hypothesis of the 2-connectivity of  $G$  by the following one: “any two vertices of  $X$  are joined by two disjoint paths of  $G$ .”

### Proof of Theorem 1.

We prove the Theorem in the following form.

“Let  $G$  be a graph of order  $n$ , and  $X$  a subset of  $G$  such that  $\delta_1(X) \geq 3$ . If any pair of vertices of  $X$  are joined by at least 2 vertex-disjoint paths of  $G$ , and if  $\delta_2(X) \geq n/p + 2$ , then  $X$  is covered by at most  $p-1$  cycles of  $G$ .”

The proof is by induction on  $p$ . For  $p=2$ , by Proposition 1,  $X$  is covered by one cycle.

From now  $p \geq 3$ . Let  $P$  be a path of  $G$  of maximum length with extremities in  $X$ . Let  $a$  and  $b$  be the extremities of  $P$ . We give an orientation to  $P$  from  $a$  to  $b$ .

Let  $X(a)$  (resp.  $A(a)$ ) be the set of vertices of  $X$  (resp. of  $G$ ) which precede the neighbors of  $a$  on the path.

As the minimum degree of  $X$  is at least 3, then  $|X(a)| \geq 2$ .

Let  $a'$  be the first vertex of  $X(a)$  different from  $a$ . There are 2 cases:

$\alpha$ ) either  $a' < z$  for every  $z \in N_P(a)$ ,  $z \neq a^+$ .

Let  $A(a') = \{z^+; z < a', \text{ and, } z \in N_P(a')\} \cup \{z^-, z > a' \text{ and } z \in N_P(a')\}$

$\beta$ ) or, there exists a neighbor  $z_1$ , of  $a$ , different from  $a^+$  such that  $z_1 < a'$ . We choose  $z_1$  as near as possible from  $a'$  and we set

$$A_1(a') = \{z^-; z < z_1, \text{ and, } z \in N_P(a')\} \cup \{z^+; z_1 < z < a', \text{ and } z \in N_P(a')\},$$

$$A_2(a') = \{z^-; z > a', \text{ and, } z \in N_P(a')\} \text{ and } A(a') = A_1(a') \cup A_2(a').$$

Let us remark that there is at most one vertex, the vertex  $a'$  in the intersection of  $A_1(a')$  and  $A_2(a')$ . By definition of  $a'$ , the intersection of  $A_1(a')$  and  $A(a)$  has at most one element. We deduce that

$$|A(a) \cup A(a')| \geq |N(a) \cup N(a')| - 1.$$

In the path  $P$ , let  $c' = p_i$  be the vertex, with  $i$  maximum, such that  $p_i$  is a neighbor of some vertex of  $A(a) \cup A(a')$ .

Let  $C$  be the cycle composed by the segment  $[a, c']$  and one, two, or three segments depending if  $c'$  is respectively a neighbor of  $a$ ,  $A(a)$  or  $A(a')$ .

We remark that no vertex of  $X \setminus C$  is neighbor of  $A(a) \cup A(a')$  by maximality of  $P$  and construction of  $C$ .

Let  $X' = X \setminus C$ . Consider the set  $D$  of extremities on  $C$  of the paths from  $G \setminus C$  to  $C$ . If  $D$  is of cardinality 1 or respectively 2 then in the case where  $|C \setminus (A(a) \cup A(a'))| = |D|$ , the cycle  $C$  is replaced by a vertex or respectively an edge.

In the general case, let us define the graph  $G'$  by removing the vertices of  $A(a) \cup A(a')$ , and adding edges  $uv$  whenever  $u$  and  $v$  are endpoints of a subpath of  $C$  consisting of removed vertices. This set is disjoint from  $A(a) \cup A(a')$  in the graph  $G$ . The transformation of  $C$  gives a cycle, say  $C'$ . The order  $n'$  of  $G'$  is at most  $n - n/p$ , and  $N_{G'}(X') = N_G(X)$ . Then  $\delta_2(G') \geq n'/(p-1) + 2$  and  $\delta_1 < X' > \geq 3$ .

One can verify that in the graph  $G'$  any 2 vertices of  $X'$  are joined by 2 vertex-disjoint paths.

By the induction hypothesis applied to  $X'$  in  $G'$ , we cover  $X'$  by at most  $(p-1)$  cycles, edges or vertices. •

### Proof of Proposition 2.

Consider a path with extremities in  $X$  and such that

- i)  $|P \cap X|$  is maximum
- ii)  $|P \cap (V - X)|$  is minimum.

Let  $a$  and  $b$  be the extremities of  $P$ . We do the following construction. For each vertex  $u \in X(a)$ , consider  $X_1(u)$  and  $X_2(u)$ , as in Theorem 1:  $X_1(u)$  is the set of vertices which succeed to the neighbors of  $u$  in  $P[a, u]$ ,  $X_2(u)$  is the set of those who precede the neighbors of  $u$  in  $P[u, b]$ .

If for some vertices  $v$  and  $v'$  of  $X(a)$ , the set  $X_2(v') \cap X_1(v)$  is not empty, then we define  $X(x)$  for every vertex in that intersection. We repeat this construction until no new vertex is given in the intersections of the form  $X_2(v') \cap X_1(v)$ . Let  $X''$  be the set of vertices of  $X$  so obtained. For every  $v$  in  $X''$ ,  $X(v)$  is contained in  $X''$ .

Let  $v_1 = a, \dots, v_t$  be  $t$  consecutive vertices of  $X''$ .

Then  $\sum_{i=1}^t (|A(v_i)| + |N_{G-P}(v_i)|) \geq \sum_{i=1}^t |N(v_i)| - (t-1) \geq \delta_t(X) - (t-1)$ , because the intersections of any pair of sets  $A(v_i)$  is contained in  $\{v_2, v_3, \dots, v_t\}$ .

Let  $x'$  be the last vertex of  $X''$ . Then there exists a cycle  $C$  which contains all the vertices of the segment  $[a, x']$ .

For  $p=2$ ,

1) Either  $\delta_1 \geq t$ , We construct the corresponding cycle and it remains  $n/2$  vertices, which is less than  $\delta_t$ . So it remains at most  $t-1$  vertices in  $X$ . They are covered by at most  $\lceil (t-1)/\kappa(G) \rceil$  cycles where  $\kappa(G)$  is the connectivity of the graph  $G$ .

2) Or  $\delta_1 \geq 2t$ , from the vertex  $b$  we define similarly a set  $X''$  and we take  $t$  consecutive vertices  $w_1, \dots, w_t = b$ . As the minimum degree in  $\langle X \rangle$  is at least  $2t$ , we may suppose that  $\{v_1, \dots, v_t\} \cap \{w_1, \dots, w_t\}$  is empty. We have

$$\begin{aligned} |\cup_{i=1}^t (A(v_i) \cup N_{G-P}(v_i))| &\geq n/2 + 1 \text{ in } G - \{b, b^-\}. \\ |\cup_{i=1}^t (B(w_i) \cup N_{G-P}(w_i))| &\geq n/2 + 1 \text{ in } G - \{a, a^-\}. \end{aligned}$$

So there is an intersection between these two sets. The path  $P$  is contained in a cycle. We deduce that  $X$  is covered by one cycle.

For  $p \geq 3$ , consider the graph  $G'$  we obtain by removing  $\cup_i (A(v_i))$  and adding edges  $uv$  whenever  $u$  and  $v$  are vertices of  $C$  bounding a segment of removed vertices. The order  $n'$  of  $G'$  is at most  $n(p-1)/p$ . Let  $X'$  be the set  $X - \cup_i X(v_i)$ . So  $\delta_t(X') \geq n'/(p-1) + t$ . We apply the induction hypothesis to  $G'$ .

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